Find My Sloths: Automated Comparative Analysis of How Real Enterprise Computers Keep Up with the Software Update Races

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Abstract. A software update is a critical but complicated part of software security. Its delay poses risks due to vulnerabilities and defects of software. Despite the high demand to shorten the update lag and keep the software up-to-date, software updates involve factors such as human behavior, program configurations, and system policies, adding variety in the updates of software. Investigating these factors in a real environment poses significant challenges such as the knowledge of software release schedules from the software vendors and the deployment times of programs in each user’s machine. Obtaining software release plans requires information from vendors which is not typically available to public. On the users’ side, tracking each software’s exact update installation is required to determine the accurate update delay. Currently, a scalable and systematic approach is missing to analyze these two sides’ views of a comprehensive set of software. We performed a long term system-wide study of update behavior for all software running in an enterprise by translating the operating system logs from enterprise machines into graphs of binary executable updates showing their complex, and individualized updates in the environment. Our comparative analysis locates risky machines and software with belated or dormant updates falling behind others within an enterprise without relying on any third-party or domain knowledge, providing new observations and opportunities for improvement of software updates. Our evaluation analyzes real data from 113,675 unique programs used by 774 computers over 3 years.

1 Introduction

Updating software in general and applying software patches in a more specific sense is a very crucial part of maintaining an ecosystem of computers safe [6]; although regular updates alone do not guarantee complete safety, falling behind
for sure poses security risks [46]. Research has shown a security patch might take months to create once a vulnerability is found; it even can come after the public disclosure of the vulnerability [25], and even after the vulnerability is publicly disclosed, many users may still use older versions. These delays open the chance for attackers to exploit those vulnerabilities; For large enterprises, such risks may lead to a significant financial loss, and a negative impact on their reputation [43].

In the recent Equifax breach case, the personal information of 143 million Americans were stolen [7], and the attacker exploited a known vulnerability, whose patch was available a few months before the incident. It could have been prevented if the software was updated in time. Another study [45] shows that more than 99% of exploited vulnerabilities were used by attackers more than one year after the vulnerabilities were publicly disclosed (e.g., CVE [15]). The WannaCry ransomware also shows how missed or delayed security updates can affect enterprises, as well as individuals [6].

Having all the users keep all the programs on their computer up to date at all times is ideal, but as shown in our probe as well as [18,25], we are very far from it in the real world; even in cases where an update is installed with minimal user involvement [18]. Numerous attempts have been made to quantify how up to date a computer is, mainly by focusing on a small set of programs [49].

The software update is a complicated process that involves multiple parties and decision factors to occur such as the availability of the machine or the connection to deliver the software update (e.g., a computer isolated with an air-gap is not updated), the system-wide policy to control update behavior in an enterprise, and each machine’s or software’s configuration (e.g., a user can stop updater due to its annoyance of notifications). We summarize the currently unsolved challenges to understand this problem as follows.

(1) First, understanding when each software’s update is created and released is important by setting the reference on the sender’s side to determine how long the update takes or it has been delayed in each client machine (i.e., on the receiver’s side). This is a real challenge due to the lack of standard channels. Several prior works [20,27,48] utilized this information by a connection to certain software companies [18,20]. Other works used a third party vendor, which attempts to collect this information using binaries’ properties. This method typically relies on the user (or software vendors) submission of the binary executable, and thus it might not provide such information for less popular or homebrew software. Our study of a large group of 774 machines shows that only 14.2% of installed software information is available on National Software Reference Library (NSRL) [4] and 75.3% on VirusTotal [47]. There is no systematic way to obtain this information for a comprehensive set of software for general usage.

(2) Second, knowing exactly when the released software has landed in a machine is a piece of crucial information to evaluate the update process of a program in each machine. Software vendors (e.g., Google) may estimate the deployment statistics if they use update management software that reports the installation timestamps or by the usage of the software (e.g., Google Chrome) if
it uses network and reports its version on usage. However, many programs
do not have such mechanisms implemented to evaluate the update processes.
They may not use dedicated update management software relying on users
to download and execute the installation package program manually. Also,
programs may not use the network or do not use telemetry functions. Given
numerous software programs being used in an enterprise, we observe many
programs in the shadow without automated well-designed update manage-
ment. Their updates solely depend on each user’s alertness or the enterprise
administration to initiate checks and updates. A systematic method to track
the update occurrences for a comprehensive list (ideally all) of software is
highly desired but missing.

(3) After all, there is no standard on how each software should be updated.
Therefore, software vendors perform updates with their own ways and own
schedules as we show in Sect. 3.2. Also, there are multiple reasons why soft-
ware is incapable of updating features; legacy software (designed without
updates), the programs made by a limited resource (e.g., small vendors or
indie software), and the terminated or outdated programs support. All these
symptoms illustrate the demand for a systematic study on how our current
software is doing with updates and guidelines suggested regarding how each
software should be managed and how an update management software should
be designed.

In this work, we attempt to fill these gaps by creating a systematic approach
to measure the update behavior of a comprehensive set of programs from all
machines in an enterprise starting with individual records and summing up to
show the overall patterns. This work solves the aforementioned challenges by
automatically estimating (1) the release time of software and (2) the landing
time of the software with a fine granularity of individual programs and individual
machines in an enterprise. From these data, we could estimate how much behind
the latest version each software in each machine is without relying on domain
knowledge of developers’ channels or third-party information specialized with
software analysis. We take (3) multiple observations out of the real data from a
real enterprise environment and provide suggestions on what would be desired
properties of an update management software. Also, we deep dive on the sloths,
that we refer to the individual software and machine behind their peers in the
progress of updates inside the enterprise environment. We attempt to measure
their risk in terms of the delay based on our inferred software versions that are
available for all programs in our observation.

To achieve these goals, we develop FMS, a tool that autonomously ana-
lyzes update patterns from the collected data, detects outdated programs and
machines, and produces timely notifications to administrators. Our evaluations
show FMS can infer version orders with 85% accuracy. Using FMS, we have iden-
tified more than 14,690 outdated programs from 774 computers and 2,705 more
programs engaged in risky behaviors.

FMS is an acronym of Find My Sloths, which refer to enterprise applications showing
undesirable delayed update behavior.
We make the following unique contributions in the analysis of software update behaviors in a real enterprise environment.

- Systematic study of software update behavior based on real-world enterprise data; Covering a total of 113,675 programs in multiple platforms observed in an enterprise with 248 people. This result brings new observations of real-world factors in software updates.
- We propose a method to estimate software release time and update delay with only data collected inside the enterprise, without relying on software vendors’ release notes or 3rd party (e.g., VirusTotal, NSRL). We enable the estimation of these information for all observed software to determine the update delay of all programs. We found the first appearance of software in a fairly large group of 774 machines can approximate software release time. We present the closeness of these two data in Sect. 3.3.
- Our approach estimates the update delay of an individual binary executable by subtracting the release time from the landing time of software update being tracked 24/7. This needs to be done individually for each software in each machine so that we can determine the delay of individual instances of software updates. It is enabled by tracking operating system (OS) events (e.g., system calls, Windows API) that access and execute each binary executable in all machines. This data enables a drill-down approach to determine the sloth in multiple layers: starting from the machine with a large update delay, we can nail down the identification of the slowest software in that machine with our fine grained individualized monitoring.

2 Observations on Enterprise Software Deployment

This research collects binary executable update records from 774 PCs and servers in an enterprise with 248 employees. This section summarizes observations from the data we collected over three years between Feb. 2017 and Feb. 2020.

**Observation 1: (Complexity of Update Transitions)** Software update state transitions are more complex at the client sides than the developers’ views. N versions of program binary distributions can cause up to $N^2$ possible transitions depending on the availability and activation of updates.

**Implications:** Even though developers release a linear sequence of binaries, *extra* non-linear transitions appear in the software update graph that defy the versioning sequences of development due to diverse situations at the client side.

After a particular version of the software is installed, the next installed version varies depending on various factors such as the updater configuration of a particular machine, the machine’s availability, connectivity, and the enterprise’s policy. When there is a gap between the installed version and the new version to be installed, some software allow a direct transition of a version to another version (e.g., 1.0 to 3.0) while some other software go through the applications
of all intermediate versions (e.g., 1.0–2.0–3.0). We have analyzed these behaviors of software in our observation in Sect. 4.

**Observation 2: (Unusual Update Transitions)** Rollbacks or regressions are not uncommon.  
**Implications:** *Counter-intuitive* transitions such as rolling back to older versions may occur in multiple software and machines due to various situations at the client side like compatibility and functionalities.

We have observed multiple software from major vendors such as Adobe Flash and Mozilla Firefox frequently show unusual transitions. We have analyzed this observation of counter-intuitive transitions in Sect. 4.

**Observation 3: (Various Update Deployment Time)** Updates released get installed at clients after various delays from less than 10 minutes to several years.  
**Implications:** In software updates, clients’ roles are as important as the developers’ release schedule because their configuration and choices decide *whether and when* updates get installed.

We have performed fine-grained analysis on these software delays running in an enterprise, individually measuring the delay of an individual software program in each machine. This result is presented in Sect. 4.

**Observation 4: (Various Support Period of Software Update)** Depending on products and vendors, we observe that the updates could be provided from no update in years to 1441 updates in three years. On average, we observe 6.4 updates for each product in our environment.  
**Implications:** Users should be aware of the risk of end-of-support software.

We have analyzed this observation in Sect. 4. Software gets updates from developers, but its degree highly varies depending on the vendor and products. Large software vendors tend to have more extended support (presumably due to their resources). Adobe Acrobat and Microsoft Office 2010 have been supported for 11 and 9.5 years so far, respectively. Software without support poses a serious risk as vulnerabilities are discovered over time.

**Observation 5: (A Long Term Software Usage)** Some software gets used for a long time. Multiple programs are used as an outdated version even after when the update support is no longer provided.  
**Implications:** Like any commodity, the software is used as far as it can perform its function. However, unlike hardware commodity, software gets vulnerabilities and becomes riskier to operate over time. Combined with the fourth observation, the software usage after the terminated support poses a high risk.

We have analyzed this usage period of software in Sects. 4 and 6.1.
3 Design of Find My Sloths

In this section, we present FMS’s design, which automatically tracks and analyzes software update patterns in a real enterprise environment. FMS is composed of three main components: (1) Automated tracking of software binary information, (2) Software update inference, and (3) Software risk analysis.

3.1 Automated Tracking of Software Binary Information

Package management systems provide automated installation, upgrade, configuration, and removal of computer programs. There exist various package management systems such as Linux Advanced Package Tool [1], RPM Package Manager [38], zypper [42], portage [19], yum [8], pacman [44], Home Brew [2], MacOS App Store [12], and Microsoft Store [32], that can possibly eliminate the user’s effort towards manual installs and updates.

However, prior study shows the limitations of update managers [27]. Although convenient to use, they do not provide full coverage of programs used on an everyday machine, nor do they guarantee the timely installation of updates, mainly due to inconveniences caused for the user during and after the update process. In addition, there exist software types that package managers cannot support, such as direct drop-in of a binary without an installer, custom installers that do not work with package managers, programs’ self-updates, downloads of related binaries or libraries, and local compilation of programs.

In this work, we do not aim to provide yet another automated updated system, but we develop an automated and unified method to monitor software updates by leveraging OS event monitoring techniques. OS event monitoring techniques have been widely used for security analysis [9,24] or software execution diagnosis (e.g., fault diagnosis, debugging, root cause analysis) [11,16]. Our monitoring agents utilize the OS event monitoring systems, which are currently available in mainstream operating systems. In Microsoft Windows systems, we use Event Tracing for Windows (ETW) [29]. We use the Auditd system [3] for Linux operating systems. Specifically, we monitor the events regarding process execution and file modification to binary executables (e.g., executable and library files). The following table summarizes the events that we use.

<table>
<thead>
<tr>
<th>Platform</th>
<th>Operating system events</th>
</tr>
</thead>
<tbody>
<tr>
<td>Windows (ETW)</td>
<td>WinExec, WriteFile, WriteFileEx, WriteFileGather</td>
</tr>
<tr>
<td>Linux (Audit)</td>
<td>execve, write, writev, pwrite, pwritev, pwrite64</td>
</tr>
</tbody>
</table>

We install the OS event tracking system (we call them monitoring agents) in hosts, including servers, desktops, and laptops, in the organization. Our agents cover various OSes and their versions—over 30 kinds—including Windows (e.g., Windows 7, 10) and Linux (e.g., Ubuntu, Redhat, CentOS) in our prototype.
Other operating systems such as MacOS can be also supported by our system by utilizing their OS event function. Once the events are collected, the rest of the process is agnostic to the operating system itself. Tracking these OS events provides the history of execution and modification (e.g., update, patch) of all software binaries on a computer in a heterogeneous manner.

![Architecture of Find My Sloths (FMS).](image)

The tracking system streams the collected events to the backend server as shown in Fig. 1. These events are stored in a database, which we use to analyze the data. Next, we explain how we use this information to infer software updates.

### 3.2 Software Update Inference with Binary Update Lineage Graph

**Lineage Graph Generation.** To track the binary update information for each software installed, we develop a lineage tracking algorithm that analyzes OS events collected by our tracking module. We identify file write events applied into binary executables and construct a binary update graph for each executable files. A file path for an executable file, which is executed, becomes a candidate for an update check. Whenever we observe a file write event applied this file candidate, this event is inferred as an update of the program.

Specifically, we first generate a graph to represent the lineage of binary updates, $G(V, E)$ where $V$ is a set of vertices which represent the information of program binaries, including both execution history and metadata. We call this graph a **Lineage Graph**. We use a SHA256 hash of the binary, $S$, as the identifier of each vertex. $E$ is a set of directed edges where each edge shows how a binary has been updated as $U(S_{old} \rightarrow S_{new})$. Each edge contains a timestamp of the update as well as the identifier of the computer that it was observed on.

Interestingly, we have observed unusual binary modification patterns in some applications, including Mozilla Firefox. To support the update without completely terminating the application, Firefox partially updates the image, reloads it in a running process, and updates the next portion of the image. It generates many bogus vertices and edges in the graph. To avoid this, we defined a time threshold, and if the sequence of binary updates happens within the threshold, we merge them in the graph and only keep the first image information. In this work, we use 99.9 percentile of the time between updates, 180s, as an empirical threshold.
Algorithm 1. Graph Creation

Input: oldSig, newSig, comInfo, path
V = [], E = []
for each item in dataList do
    u = null, v = null
    if (oldSig ∈ V) then
        u = V.getItem(oldSig)
        add comInfo and path to u
    else
        u = new Node(oldSig); V.add(u)
    if (newSig ∈ V) then
        v = V.getItem(newSig)
        add comInfo and path to v
    else
        v = new Node(newSig); V.add(v)
    if ((u, v) ∈ E) then
        add comInfo and date to e
    else
        E.add (new Edge(u, v))

Global Lineage Graphs as an Enterprise Level Update Summary. We then generate the graphs for all hosts in the central server to create a collective graph and comparatively measure the risk factors for each host. We use a binary signature, S, to identify the same binary in different hosts and construct a global update graph for each application to visualize binary update patterns across hosts. It allows us to have a bird’s eye view of how each program is updated, and similarities and differences of update patterns across hosts. Algorithm 1 presents how we construct global graphs.

Figure 2 demonstrates a simplified example of how we construct a global view of update patterns. In this scenario, we have four hosts that have the execution or modification history of a particular application. The application has three distinct binary signatures, the version 64.0 (S_{64.0})^2, S_{65.0}, and S_{65.0.1}. Host 1 has execution and update history where the binary has updated at t_1 from S_{64.0} to S_{65.0} and updated to S_{65.0.1} at timestamp t_2. The binary in Host 2 has updated to S_{65.0} at t_1 but never updated to S_{65.0.1}. Host 3 only has an execution history of S_{65.0}, and Host 4 does not have any sign of S_{65.0}, but the binary has directly updated from S_{64.0} to S_{65.0.1} at t_2. Note that this is simplified graphs to demonstrate, but original nodes contain metadata and execution history (e.g., timestamps when the binary has executed).

We collect graphs from each host and construct a global graph aggregating update patterns from all hosts. The global graph shows how the binary has updated, who have out-dated binary executables, and how frequently out-dated binaries have executed (we elide an execution history). We discuss how we measure the risk factors for each host using the global view in the next section.

Tracking Binary and Library Update. We detect updates based on overwriting of the binary files. These can be either by a write event or an execution.

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^2 This version number 64.0 is presented only for an illustration purpose. A lineage graph is constructed using binary hashes and their appearance orders without using the software’s specific version numbers, which may not always available or accurate.
event which shows a new signature for the executable image. Note that not every program update changes the main executable. It is possible to update library files while the main executable stays unchanged. Most Windows applications developed by .NET framework update the manifest in the main executable even if the update is only applied to a library file [30,31], and thus FMS can detect it. However, if the update is solely through the dynamic library without updating the manifest, or main program FMS cannot detect it, and we leave it as our future work. We further discuss this issue in Sect. 9.

Table 1. NSRL Program and Version Coverage. Latest versions available for four popular programs as of the date of the writing and their corresponding records in the latest NSRL [4] release (RDS Version 2.71 - Dec. 2020) showing NSRL release is not up-to-date for several actively developed software programs.

<table>
<thead>
<tr>
<th>Program name</th>
<th>Latest</th>
<th>NSRL versions</th>
<th>Program name</th>
<th>Latest</th>
<th>NSRL versions</th>
</tr>
</thead>
<tbody>
<tr>
<td>Mozilla FireFox</td>
<td>73.0</td>
<td>72.0, 68.0, 61.0.1</td>
<td>Acrobat Reader DC</td>
<td>21.001.20138</td>
<td>N/A</td>
</tr>
<tr>
<td>Google Chrome</td>
<td>80.0</td>
<td>76.0, 49.0, 47.0</td>
<td>Sublime</td>
<td>3.2.2</td>
<td>N/A</td>
</tr>
</tbody>
</table>

Limited Reference Services. We have considered using reference libraries for software such as National Software Reference Library (NSRL) [4] or VirusTotal File Search [47] that can retrieve information of application by a hash value. They often provide known-vulnerabilities (e.g., CVE [15]) and it can be directly used for the risk prediction. Table 1 shows version information for 4 popular programs in NSRL. As shown, these references can be used with some mainstream applications, however, they often fail to provide information on less-known applications or free software distributed via source code. Also the fact that less popular programs have less disclosed vulnerabilities [5] does not mean they are more secure and they still should be monitored. We submitted all executable files installed in our enterprise to NSRL and VirusTotal database, and we found that only 14.2% of installed binary information is provided by NSRL and 75.3% of binary information is available in VirusTotal.

3.3 Software Risk Inference

In order to estimate the risk each computer or software imposes on the enterprise environment, the state of software on individual computers, as well as, the overall enterprise status as a whole have to be quantified with a score; we introduce a metric based on the delay in update time, that is defined as the estimated delay of the installation time behind the latest version in its software distribution.

FMS monitors the process and file events and creates lineage graphs as introduced in Sect. 3.2. From the data stream, we catch new updates as they emerges. These updates are correlated and aggregated over all the enterprise to construct the global update graph of a program. To give the administrators a better perspective, we also provide them with information on how common each program
is in the enterprise and the update patterns for that particular program. Here is how FMS estimates the risk caused by the software update delay.

**Inferred Versions.** Relying on software vendor’s release schedules or third-party services’ information does not scale to cover the majority of software. Therefore, our approach determines the risk estimated by only using the measured metrics in our environment. In this scheme, the versions are determined by *temporal appearance order of signatures* relative to the prior edge nodes in an enterprise. We found that the statistical data collected from 774 machines approximate the actual order of versions with 85% accuracy. We calculate accuracy by counting the number of versions found in the correct order by FMS as compared to their respective compilation timestamp in a test set of programs related to 16,273 updates chosen randomly from Windows binary files.

Let us define a program *p* has a set of its signatures \( V_p = \{S_{p,v,m}\} \) where \( v \) is a version and \( m \) is a host index in an organization. Among multiple signatures observed, we can find the one leading the version transitions that appear most recently compared to other versions and transitioned from one or more edge nodes of the graph. We call this version the head \( \hat{v} \), and its signature is represented as \( S_{p,\hat{v},m'} \) where \( m' \) is the host having this version. This signature is determined from the graph structure.

**Estimated Release Time, Head Arrival Time, and Arrival Time of Updates.** To estimate the update delay, we need to determine three types of timestamps. The first one is the release time, \( r(S_{p,v,m}) \), when the software vendor begins to distribute a new version of the software. This information is hard to obtain with comprehensive coverage of the majority of software. We estimate this information with the first appearance time of a signature, \( S_{p,v,m} \), in an enterprise. The second metric is the deployment time of the latest estimated version, \( d(S_{p,\hat{v},m'}) \), that is the installation time of the most desirable version where the function \( d(s) \) gives the installation time of a signature \( s \). This is determined after the head signature is found from the topology of the graph. The third timestamp is the actual installation time, \( d(S_{p,v,m}) \), that is a specific time stamp when the software binary signature, \( S_{p,v,m} \) is installed. Its absolute installation time since the release is \( d(S_{p,v,m}) - r(S_{p,v,m}) \).

**Update Delay Metric.** The update delay of software *p* of the current version \( v \) in a machine \( m \) behind the head \( \hat{v} \) is defined as the following formula.

\[
D_{p,v,m} = |d(S_{p,\hat{v},m'}) - d(S_{p,v,m})|
\]

This value measures a relative delay compared to the installation of the latest version within an enterprise. Since this is calculated only by using our measurements, we can generate this risk score without any domain knowledge requirement, such as the release date and time, which is only available by developers. Such wide applicability to all programs is a unique strength of our approach. We present interesting comparative statistics showing how each software in a particular machine is doing compared to all other software instances within an organization in Sect. 4. To this end, FMS generates a daily report that outlines
the state of the risk across the enterprise and the list of machines that contain outdated programs with detailed risk information.

4 Characteristics of Software Updates

Data Set. We collect binary access history (i.e., execution and modification history) from a Fortune 500 Tech company. The company has 248 full-time employees in multiple divisions, including research and technological development (RTD), maintenance, financial, and human resources. We have installed our binary monitoring module (Sect. 3.1) on 774 computers comprising 591 Microsoft Windows and 183 Linux machines to monitor binary execution and modifications over three years. In total, we observed 113,675 unique programs with an average of 305 observed on a computer and a total of 40,971 updates.

Binary Distribution. To better characterize the distribution of programs on computers, Fig. 3 shows a histogram for the number of unique programs in each computer. There is an average of 305 programs executed on each computer in the period of our observation.

Figure 4 shows the number of binary executables published by different publishers. We leverage the metadata from the Windows binaries to identify the publisher for each binary. We only count images that contain publisher information. This shows the imbalance between popular companies and companies with less popularity in our sample enterprise. This figure shows a sample of the meta information shown on the system’s user interface. It also shows how diverse vendors are and how each one’s dominance differs from the other. In this graph Microsoft followed by Adobe and Google have the highest number of programs.

Binary Updates. We have observed 40,971 updates of binary executables from 774 Windows and Linux computers we monitored during this study. These updates collectively introduced 11,948 new binaries to the enterprise. Updates once released are not always installed within the same period on all computers, Fig. 5 shows how each update is installed on different computers.

Fig. 3. Histogram over the number of computers for each unique program appearing.

Fig. 4. Number of unique software products and unique versions observed for each vendor, each dot represents one vendor.
Fig. 5. CDF for the time over which each update has been applied to different computers. Multiple products are shown along with the overall list.

Fig. 6. Percentage of in-order and out-of-order updates for 5 different products. The number in the parentheses shows the number of updates.

Update Order Patterns. We call any update that does not follow the release order an out of order update, e.g., going from 1.0 to 3.0, skipping a known intermediate version 2.0. Furthermore, as mentioned in Sect. 2, there are a noticeable amount of downgrades observed in our dataset as well. We consider them risky behavior. Although they might have happened for reasons like a downgrade after a feature has been removed from the new version or user’s choice, they could pose a risk to the enterprise. Figure 6 shows how common these out of order updates and downgrades are among some well known products. The programs in this figure observe a total of 55 downgrades. They took programs to versions with 19.7 CVEs on average at the time of downgrade.

Product Lifecycle. We study the lifecycle of programs in our observation period by their compilation timestamps which show how spread the versions we have observed are. Figure 7 characterizes each program’s behaviors in terms of the number of days in which we have observed updates and the number of unique versions we have observed in our data collection. The highlighted nodes represent (0) Microsoft Malware Protection, (1–3) executable binaries included in Microsoft Windows Operating System.
Another part of the lifecycle is the delivery of the product to users. We observe the delay between the programs’ installations and their releases in our dataset varies significantly by vendor and product. Figure 8 shows some famous vendors and their delivery performance. For example, Google has a short delay updating 89% of installations within 18 days while Mozilla has a long deployment time where some versions got installed after several months. This may have to do with the popularity of the software.

**Update Propagation.** Updates are not installed on all computers at the same speed. The speed at which the update is installed across computers in the network depends on multiple factors such as the application, the user, configurations, and the vendor’s delivery method. Figure 5 shows how the updates of four programs are installed compare to the overall distribution. We can see ‘Microsoft Malware Protection’ updates are distributed faster on the computers compared to ‘Mozilla Firefox’. The vertical bars of each color represents the time when the maximum is reached.

### 5 Evaluation of Software Update Risk

This section presents the estimated risk determined by the delay in updates compared to peers in an enterprise environment. Figure 9 presents the estimated software update risk in our environment with three scopes from the broad view across all hosts (top) to the programs in a specific machine (bottom two).
Fig. 9. Risk analysis on computers (top) shows average delays (dots) and the short horizontal bars show the maximum delay for each computer. Expanded views show how programs on a computer stand against the enterprise distribution (numbers show total computers the program has been observed on). Plots marked by 1,2 show the details of computers marked with corresponding numbers on the top plot.

Enterprise-Level View. The top figure in Fig. 9 shows the risk scores estimated for all host machines, which show an aggregated risk score for each machine. This view highlights the uniqueness of our approach making the risk score available for all observed machines because our approach is agnostic to OS and programs. Each vertical bar (I) shows the maximum and minimum risk scores as the top and bottom marks and the average risk score of all software within each machine as a dot in the middle. We can observe many machines have one or more highly risky programs indicated by the maximum value far away from the average, which serves a comparable score across peers. We placed a blue box to show where the averages of most machines stay. We select one sample standing out with high risk shown as a higher average (noted with the red 1 sign), and another sample of a lower risk score which stays together with peers within a blue box (noted with the red 2 sign). We provide drill-down views of these two machines with a comparison next. These two particular machines are chosen to illustrate a contrast with a comparable set of programs.

Machine-Level View. The next two figures in Fig. 9 present the views of particular hosts regarding all software programs running in them. Again, this highlights our unique view of a comprehensive coverage of all programs’ risk scores
because we do not rely on developers’ or third-party’s program metadata (e.g., reputation) to determine the risk score. Each bar represents the update status of one software in all machines. The total number of machines having this software executed is written on top of the bar. Each machine has a score of delay showing how much its version is behind the head (the latest version in an enterprise). First, the average of all delays is calculated. The number of machines having the delay shorter than the average is shown as an empty bar, and the machine having the delay longer than the average is shown as a bar with the hatches. The red horizontal lines indicate where this machine is among multiple peers. The first graph (note 1) in the middle indicates multiple red lines are placed inside the hatch bars meaning that such software has its version more than the average delay among peers. For instance, multiple programs such as Adobe Flash player, AMD External Events, Microsoft Edge, Microsoft Outlook, and .NET framework show high risks. Administrators can use this view to spot vulnerable software. The second graph (note 2) in comparison shows most software programs are up-to-date by having the red horizontal lines at the bottom.

6 Case Study

This section presents an instance of desirable software management and an undesirable example of a software product. These examples show how different vendors’ software management policies are, highlighting what needs to be done.

6.1 An Example of a Desirable Software Management

In Sects. 4 and 5, we presented our observation on numerous shortcomings in how vendors design the software distribution policies and actions to protect their customers. We present Google Chrome as a desirable example with the speedy delivery of software ensuring that most of the computers are updated to the latest version in the shortest time.

Figure 10 illustrates how 75 different versions of Google Chrome have been updated. The top figure shows the deployment of each version over time. While

Fig. 10. Software update management of Google Chrome browser updating 75 Google Chrome versions over the period of 35 months. The top plot shows the version transitions via the deployment of each version of Chrome with respect to the number of computers. The bottom plot shows the number of CVEs discovered for each version.
several versions have lingered around over the maximum period of 10 months, a package was replaced by a next version after less than a day in the minimum case and after 26 days on average.

The bottom figure shows the number of CVEs accumulated for each version over time. For this case study, we manually collect reported CVEs between April 2017 and April 2019 for Google Chrome and cross-match them with versions we observed in our dataset. The total number of CVEs grows to more than five hundred for version 52.0. This data illustrates that the only solution to operate software safely is to update it with the version whose vulnerabilities are not yet discovered for the time being (although it would end up with a similar situation as the prior version). While more and more vulnerabilities have been disclosed, we can observe Google has been proactively producing new versions that replaced most of the vulnerable installations promptly.

6.2 An Example Undesirable Software Management

An opposite situation happens in another software called Viber. Viber is a communication app that has a desktop version. Unlike the previous case, this software vendor does not follow up with the management of software. In fact, this software is not shipped with its own updater. Therefore, a user will need to download and install the updated version by herself manually. This is a very ineffective way to manage software because of multiple reasons. First, most users are not likely to check new versions of software (e.g., by visiting the software website, etc.) with concerns about vulnerabilities. Second, even though the users may become aware of a new version (e.g., registration and a newsletter), they may not be capable of downloading a newer version and installing it.

We observe Viber installed on 4 machines, and they have been consistently executed during our observation period. Of these installations, only one machine was updated twice until Jan 2018, which leaves the machine out-dated for more than 2 years in which the application has been used regularly. During this time, 2 CVE have been disclosed that directly affected Viber Desktop and 4 more that indirectly affect the Viber platform.

Some vendors seem to improve behaviors with time, as shown in Fig. 11 for instance, Skype begins with versions not being updated and some versions lingering on for a long time in 2017, but in 2019 turns to a desirable update behavior where a majority of computers receive their updates in time. On the other hand, however, Firefox, also shown in Fig. 11 maintains bad update habits to the extent where in no time more than 10 of our sampled machines have the same version installed. Microsoft Edge has the worst performance among observed browsers. Edge versions maintain their dominant presence for more than 10 months and we observe major activity from them on these machines.

7 Lessons Learned and Our Suggestions

This section summarizes what we learned from observations and analysis of the real-world enterprise data, and we made a few suggestions for better security.
Minimizing Update Deployment Time. One important issue in updates is deployment time. In the end-to-end point of view, collective effort is necessary to minimize the deployment time by multiple parties including developers, system administrators, and end users.

Reliable Update Transitions. Given $N$ versions ($s_n \in S = \{s_1, ..., s_N\}$), $N^2$ transition patterns may occur depending on current and next versions. When there is a gap, a direct transition to the latest can ease deployment regarding compatibility and time in subsequent updates if the update is stand-alone. If it has to be incremental, developers should thoroughly test every transition.

Software Downgrading. We have observed a considerable number of version downgrading in our data set. This is also in line with findings of [27] that users partially steer away from updates due to the issues with compatibility or features.

Automated, Enforced, Silent Delivery. We observe the effectiveness of silent delivery that the user has no interaction with the download and installation of the update. Google Chrome shows a successful practice of silent delivery.

Retirement Plan and the End of Support Notice. Proper announcement and communication of product support (e.g., a good example is the Ubuntu’s LTS plan) is important to know the software’s managed period. We believe that an automated method to query a program’s lifespan would help the user avoid end-of-support programs.

8 Related Work

Risk Prediction and Vulnerability Assessment. RiskTeller [13] creates computer profiles based on file appearances and the general use of the machine and predicts each machine’s risk based on its installed software and the usage history. RiskTeller uses external information, including VirusTotal and CVSS
Table 2. Comparison of FMS and related work. Cov∀: Applicable to all programs, No D.K.: No third-party domain knowledge (e.g., VirusTotal, Symantec) required, Applicability: the data can be generated from an enterprise and the method is applicable.

<table>
<thead>
<tr>
<th>Name</th>
<th>Prediction</th>
<th>Cov∀</th>
<th>No D.K.</th>
<th>Data period</th>
<th>Applicability</th>
</tr>
</thead>
<tbody>
<tr>
<td>Xiao et al. [49]</td>
<td>Vul. exploitation</td>
<td>✗</td>
<td>✗</td>
<td>1.5 years</td>
<td>✗ (Botnet, patches)</td>
</tr>
<tr>
<td>Sharif et al. [39]</td>
<td>Security incident</td>
<td>✗</td>
<td>✗</td>
<td>3 months</td>
<td>✗ (Cellular provider)</td>
</tr>
<tr>
<td>Ovelgönne et al. [35]</td>
<td>Infection</td>
<td>✗</td>
<td>✗</td>
<td>8 months</td>
<td>✗ (Symantec data)</td>
</tr>
<tr>
<td>Tiresias [40]</td>
<td>Security event</td>
<td>✗</td>
<td>✗</td>
<td>27 days</td>
<td>✗ (Symantec data)</td>
</tr>
<tr>
<td>ASSERT [34]</td>
<td>Intrusion</td>
<td>✗</td>
<td>✓</td>
<td>Test data</td>
<td>✗ (Test data)</td>
</tr>
<tr>
<td>Liu et al. [26]</td>
<td>Security incident</td>
<td>✗</td>
<td>✓</td>
<td>2 years</td>
<td>✗ (Network data sets)</td>
</tr>
<tr>
<td>Kang et al. [22]</td>
<td>Malware infection</td>
<td>✗</td>
<td>✗</td>
<td>640 days</td>
<td>✗ (Symantec WINE)</td>
</tr>
<tr>
<td>RiskWriter [10]</td>
<td>Security posture</td>
<td>✗</td>
<td>✗</td>
<td>1 year</td>
<td>✗ (Symantec data)</td>
</tr>
<tr>
<td>Find My Sloths</td>
<td>Vul. program risk</td>
<td>✓</td>
<td>✓</td>
<td>3 years</td>
<td>✓ (General OS log)</td>
</tr>
</tbody>
</table>

score [15], to enable an accurate prediction. Other studies infer the risk of a computer (or mobile device) by using user behaviors of web sessions [39], distribution of binary files on the computer, the user’s mobility between different ISPs [35], modeling the security-related decision-making process of users [37], and permission and internal activities of Android applications [41,50].

RiskWriter [10] and Liu et al. [26] measure the risk for an enterprise based on its externally measurable metrics such as DNS misconfigurations. Kang et al. [22] predict malware infection epidemics based on antivirus software telemetry information. Another direction of risk prediction [34,40] is to predict future security incidents based on the history of the past security-related events.

Kotzias et al. [23] studies patching landscape for the enterprise; however, the set of programs they tracks is limited to 124 well-known server and host programs for which they obtain hash and version information for the main executable. Nappa et al. [33], also studied effects of patching, threats of multiple installations of a software and the shared libraries from 10 popular programs at scale.

Unlike previous studies (Table 2), we infer update delay using only data collected inside the enterprise, without relying on external information, such as software vendors’ release notes, third party metadata of software, or known vulnerabilities.

**Update Development.** There exist works that focus on the development side of updates that study how well developers respond to the needs for updates, including security vulnerability disclosure [25]; or the response speed of developers to new requirements of ever-changing platforms. Code change analysis [28], commit message analysis [28,36], vulnerability disclosure analysis [36,49], and versioning analysis [14] techniques have been used in these studies.

**Update Deployment.** Software development companies often have their solutions to deploy updates and patches effectively and promptly. For instance, Microsoft’s On-Premises Update Server provides a unified interface, and it tries
to minimize the network traffic by caching the update [20]. Google applies silent and automatic updates [17] to their products, including the Google Chrome web-browser. There exist studies [27, 48] that evaluate the effectiveness of software update deployment models. They use a mix of user studies and monitoring a small set of programs to find the challenges stopping an effective update process. A recent study [21] identifies design flaws or malicious installations in the automated software deployment systems.

9 Discussion and Future Work

This section discusses the limitations of our work and possible improvements.

Reasons Behind Update Behaviors. In this study, we focus on identifying software update behaviors and estimating risk; however, understanding why a certain update behavior happens is a remaining challenge beyond our scope. For instance, we detect software downgrade or rollback from the lineage graph, but the reason why that decision has been made is unknown. In the future, we will conduct a comprehensive user study and a deep analysis of update management platforms to understand the reasons behind unusual update behaviors.

Grouping Updates Installed in Varied Paths. The current version of FMS detects an update of software individually based on the path where the update is installed. However, we rarely observe software that installs the update in a separate path while keeping the previous versions of the binary. It then modifies a symbolic link file to point the binary to execute. Currently, we may not correctly detect such update behaviors. We plan to implement a merging mechanism to automatically find updates in separate paths to address this limitation.

Updates By Library Modification. The current FMS cannot detect updates that only modify the library but not the main executable. We plan to improve FMS’s monitoring module to enable tracking library import information at runtime. Then we will track all updates to those library files.

Leveraging Security Information. We plan to optionally incorporate standardized security information, such as code signing, to further enhance the risk estimation in FMS for certain programs whose metadata are available.

10 Conclusion

Software lifecycle management is a complex and costly process that needs the dedication of a software vendor. When it is not properly done, we find that it directly leads to risk for every software as vulnerabilities are discovered over time. We propose an automated approach to analyze the entire software updates in an organization comprehensively achieved by utilizing only observed metrics instead of relying on developers’ or third-party software release information and metadata. Our evaluations shed light on the current industry practices on how they manage software due to high coverage of update risk assessment on total
113,675 programs’ 40,971 updates in 774 machines used by real people daily. We organized our comprehensive evaluations on the current software updates practices, which suggest a list of desired design decisions on update management software to operate software securely.

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